

# New Yoyo Tricks with AES-based Permutations

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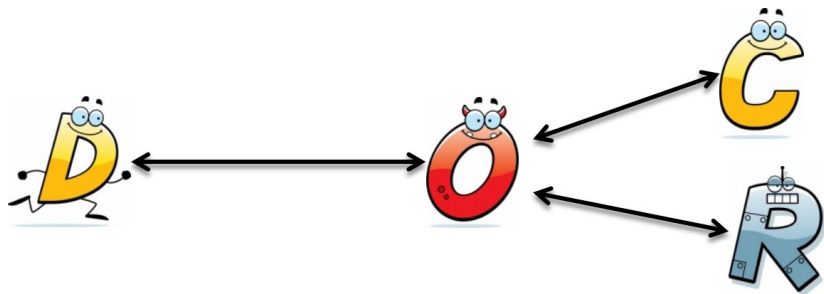
# The Problem: Devising Distinguishers

Distinguish between **what** and **why**?

Exhibiting Non-random Behavior

# The Distinguishing Setting

1. D tries to distinguish between C and R
2. Can make queries to O
3. O behaves as either C or R
4. At the end D has to guess who is O impersonating
5. D wins if its guess is right



# Lets play a Game

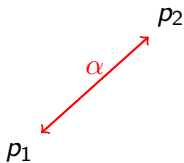
**Setting:** Adaptive Chosen Plaintext/Ciphertext

Will look similar to Boomerang Attack



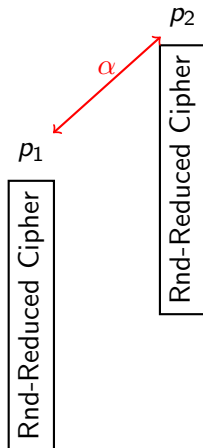
Select messages  $p_1, p_2$  with  $p_1 \oplus p_2 = \alpha$

Is there a special way to choose  $\alpha$ ?



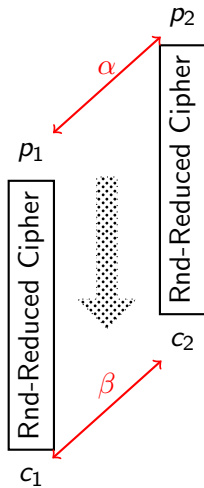
# Apply some rounds of some cipher

How many rounds? What type of cipher?



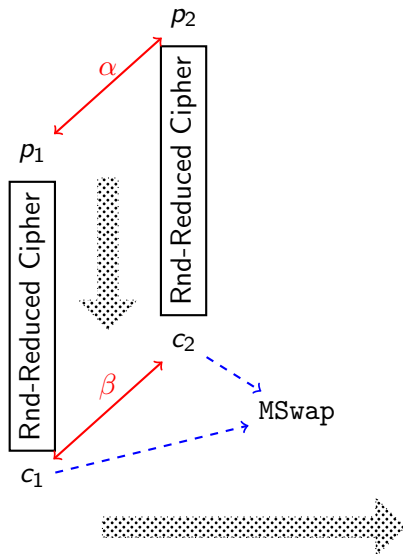
Get  $c_1, c_2$  with  $c_1 \oplus c_2 = \beta$

$\beta$  is the ciphertext difference



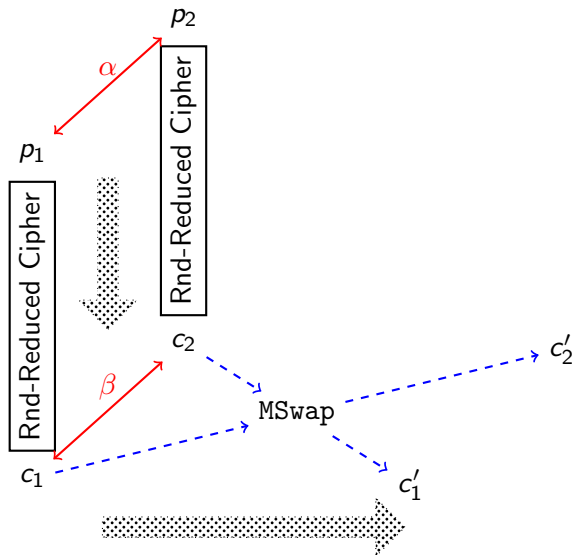
# Use MSwap to swap bytes/words of $c_1, c_2$

How does this swap work?



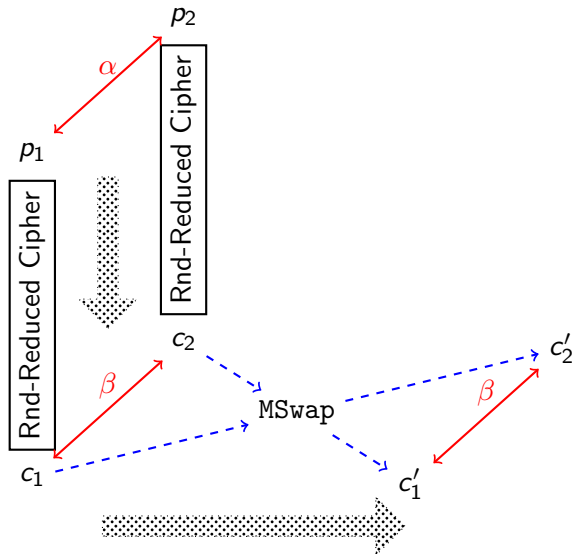
# Generate new ciphertext pair $c'_1, c'_2$

What is the relation between  $c'_1, c'_2$ ?



Invariant:  $c'_1 \oplus c'_2 = \beta$

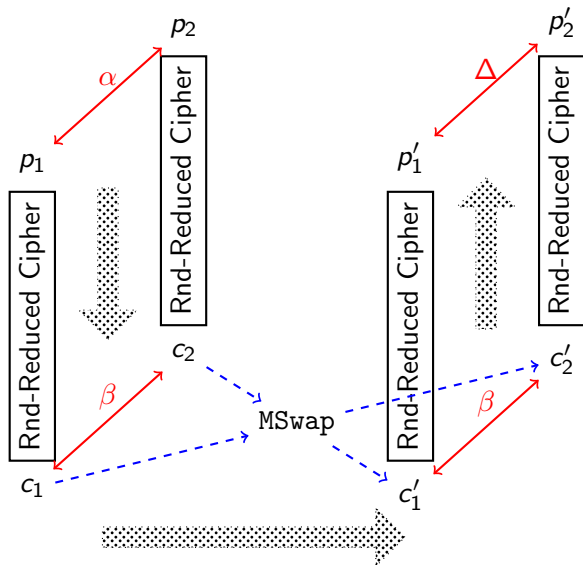
How does this part differ from the **Boomerang** Attack?





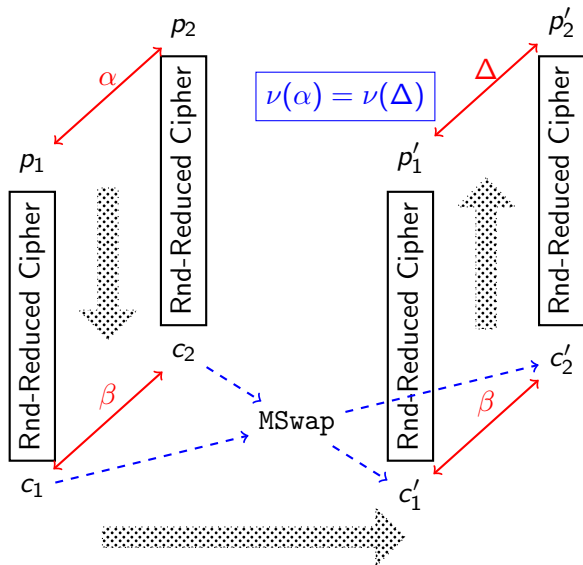
Get  $p'_1, p'_2$  with  $p'_1 \oplus p'_2 = \Delta$

Does  $\Delta$  have a special property?



# Hypothesis: Property $\nu$ induced in $\alpha$ is preserved by $\Delta$

What is this **property**  $\nu$ ?

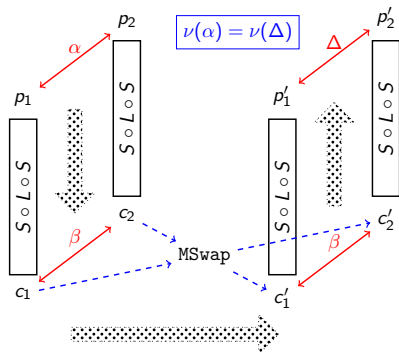


- ▶ Is there a special way to choose  $\alpha$ ?
  - ▶ Zero Difference Pattern (ZDP). Explained later.
- ▶ How many rounds? What type of cipher?
  - ▶ **2-Rnd Generic SPN**
- ▶ How does the swap work?
  - ▶ Swap based on non-linear layer. Explained later.
- ▶ Does  $\Delta$  have a special property?
  - ▶ Same as  $\alpha$
- ▶ What is this **property**  $\nu$ ?
  - ▶ Zero Difference Pattern (ZDP)

$$G'_2 = L \circ S \circ L \circ S$$

Two full generic Rounds

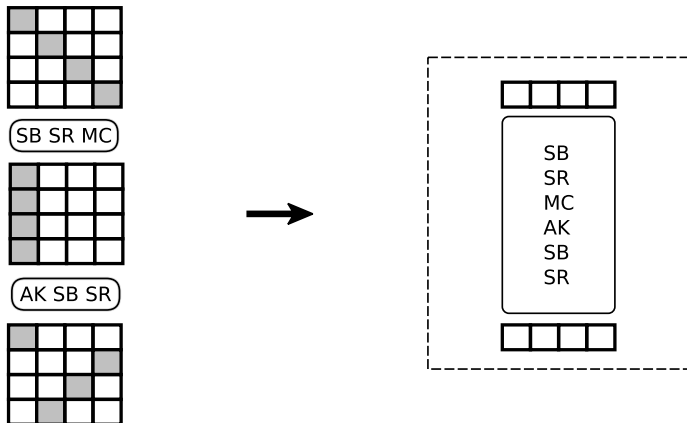
$$G_2 = S \circ L \circ S \quad \leftarrow \text{Dropping final linear layer (to simplify)}$$

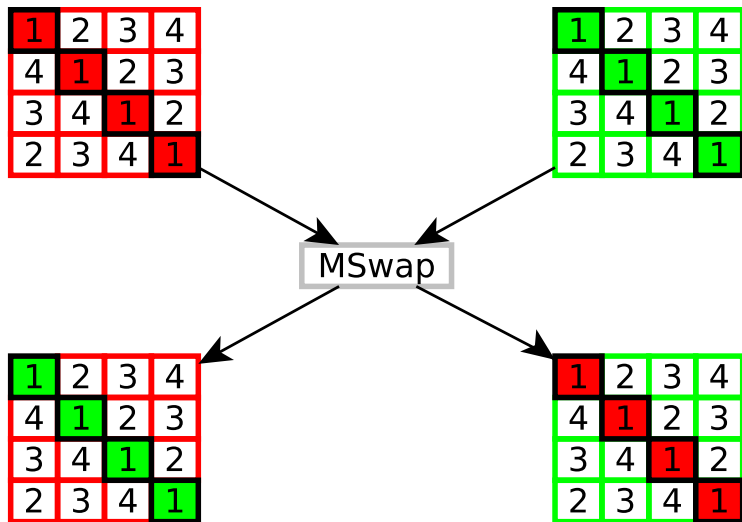


- ▶ ZDP of  $\alpha$  is preserved by  $\Delta$

Applied to AES

- ▶ First Key-independent Yoyo Distinguishers of AES
- ▶ 5-round practical key recovery





$$p_1 = \begin{array}{cccc} \text{fa} & \text{b1} & \text{5a} & \text{2f} \\ \text{b7} & \text{64} & \text{0e} & \text{f1} \\ \text{f8} & \text{9f} & \text{22} & \text{15} \\ \text{28} & \text{87} & \text{32} & \text{25} \end{array}$$

$$p_2 = \begin{array}{cccc} \text{2e} & \text{b1} & \text{5a} & \text{2f} \\ \text{b7} & \text{70} & \text{0e} & \text{f1} \\ \text{f8} & \text{9f} & \text{f2} & \text{15} \\ \text{28} & \text{87} & \text{32} & \text{4c} \end{array}$$

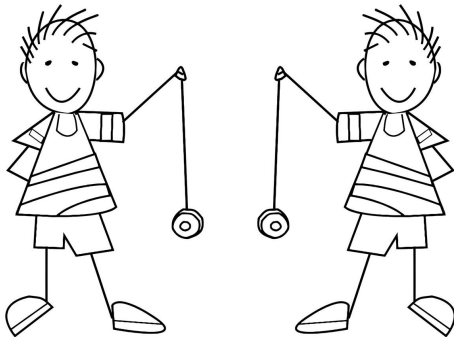
$$\alpha = p_1 \oplus p_2 = \begin{array}{cccc} & \text{d4} & 00 & 00 & 00 \\ & 00 & \text{14} & 00 & 00 \\ & 00 & 00 & \text{d0} & 00 \\ & 00 & 00 & 00 & \text{69} \end{array}$$

$$ZDP(\alpha) = \{0, 1, 1, 1\} \quad wt(ZDP(\alpha)) = 3$$

- ▶ New pairs of plaintexts and ciphertexts are made **adaptively** from the original pairs.
- ▶ While making new pairs a **certain property** is kept invariant.
- ▶ A common strategy is the use of **zero difference** in the pairs.
- ▶ An invariant property is verified at the end



# Our Aim: How To Exploit Yoyo Further



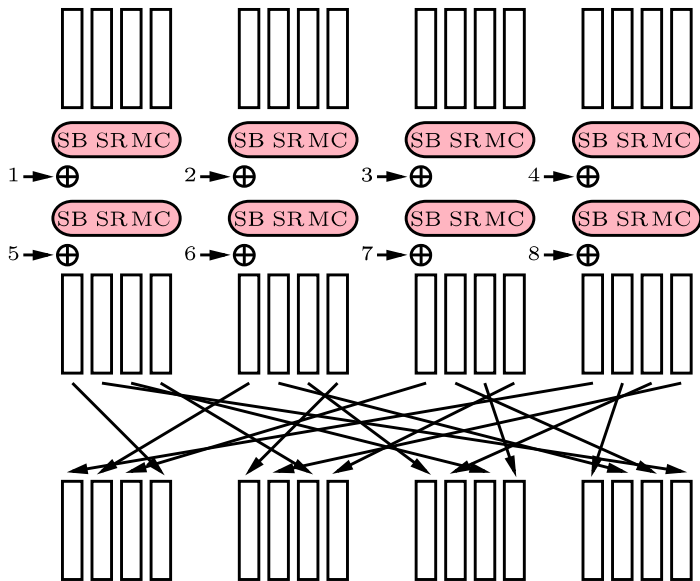
Our Target: AES-based Public Permutations

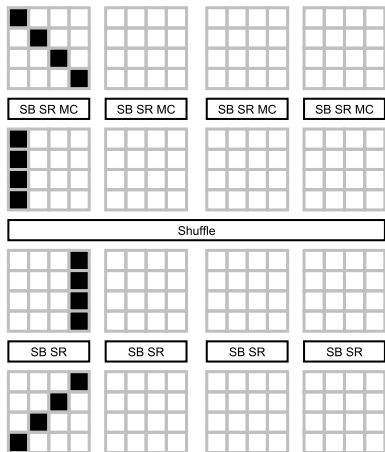
## AESQ Permutation

- ▶ Internal permutation of AE scheme PAEQ
- ▶ PAEQ  $\leftarrow$  2nd Round CAESAR candidate
- ▶ By Birukov and Kovratovich

## AES in *Known-Key* Setting

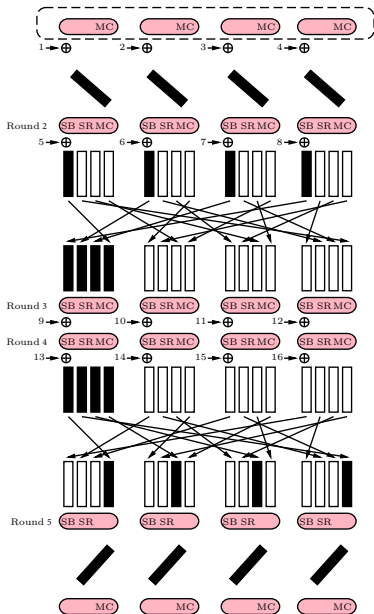
- ▶ Known-key paradigm
- ▶ By Knudsen and Rijmen
- ▶ Under Known-key AES behaves as a public permutation





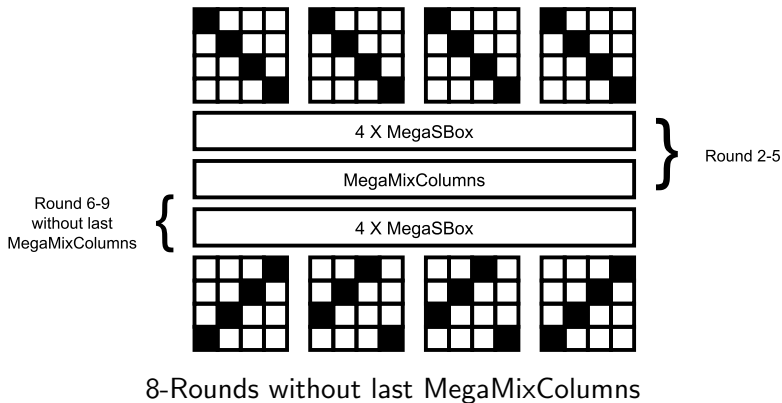
## 32-bit SuperSbox

- ▶ 16 SuperSBox-es
- ▶ Cover 1.5 Rounds
- ▶ **Must start from even round**



## 128-bit MegaSBox

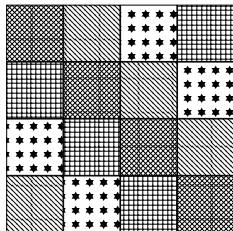
- ▶ 4 MegaSBox-es
- ▶ Cover 3.5 Rounds
- ▶ **Must start from even round**



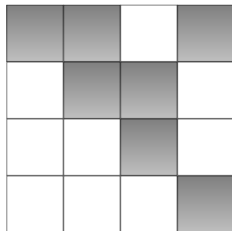
$\nu^2$ Introducing **Nested** Zero Difference Pattern $\alpha \leftarrow$  Sample State

$$\nu(\alpha) = (0, 0, 1, 0)$$

$$wt(\nu(\alpha)) = 1$$



A sample state

 $\alpha_0$  $\alpha_1$  $\alpha_2$  $\alpha_3$ 

Active Byte



Inactive Byte

$$\nu_1^2(\alpha_0) = (0, 0, 0, 0),$$

$$\nu_2^2(\alpha_1) = (0, 0, 1, 1),$$

$$\nu_3^2(\alpha_2) = (1, 1, 1, 1),$$

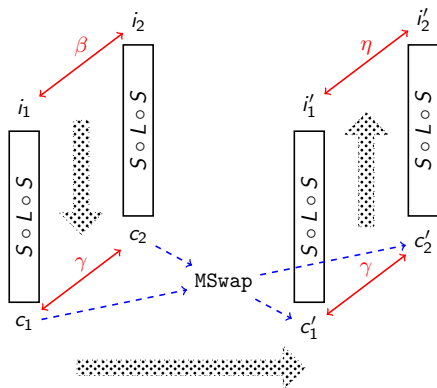
$$\nu_4^2(\alpha_3) = (0, 1, 1, 1)$$

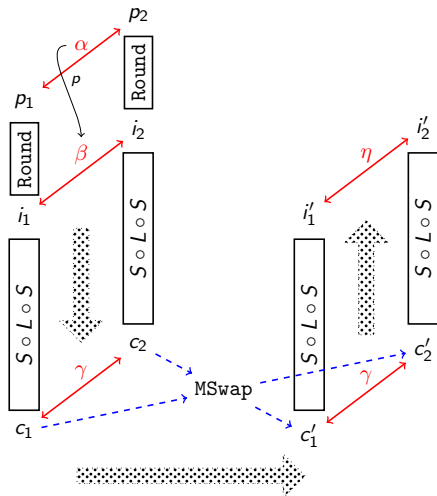
$$wt(\nu^2(\alpha)) = 9$$

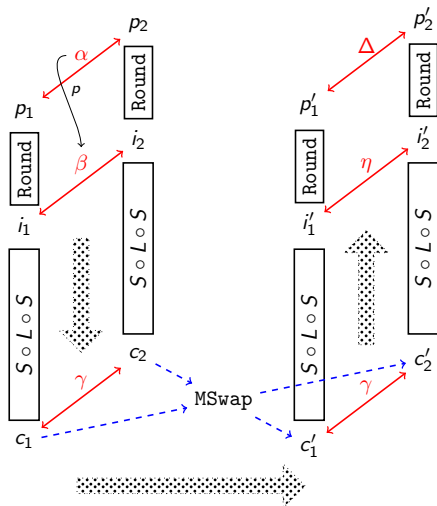
# Strategy 1: Prepend-Append

Probabilistic Yoyo

Using Classical Differentials

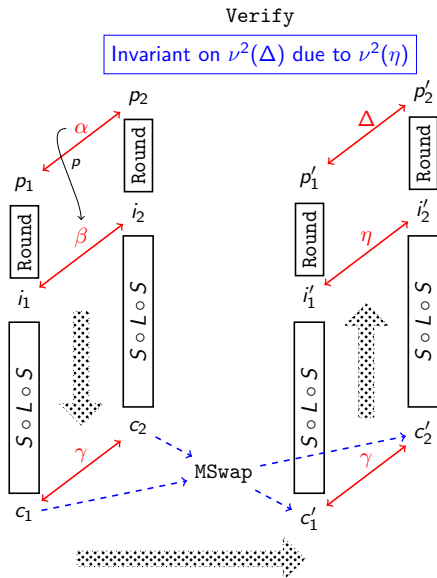






# Probabilistic Yoyo Distinguisher

Property verified on  $\Delta$

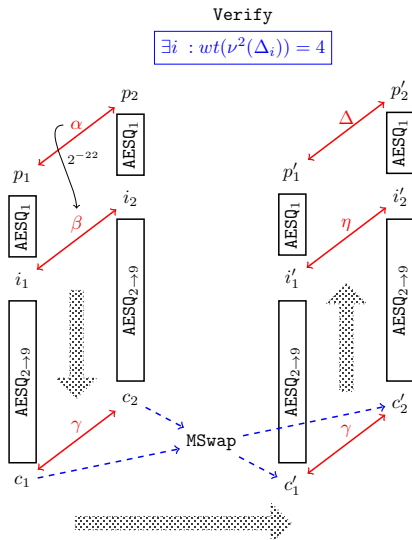
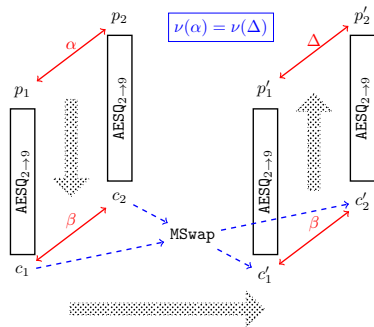


# Application: AESQ

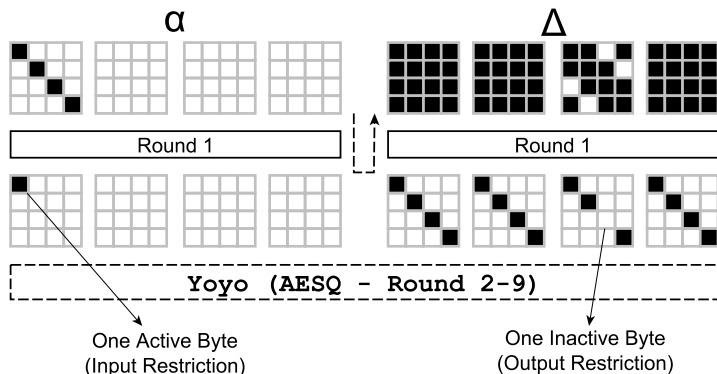
First 9-round Distinguisher starting from Round-1

Practical Complexity

## Basic Yoyo 8-Rounds



## 1-Round Extension



For  $\text{AESQ}_{1-9}$

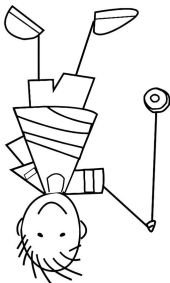
$$\Pr[\exists i : wt(\nu^2(\Delta_i)) = 4] = 2^{-26}$$

For  $\mathcal{R}$

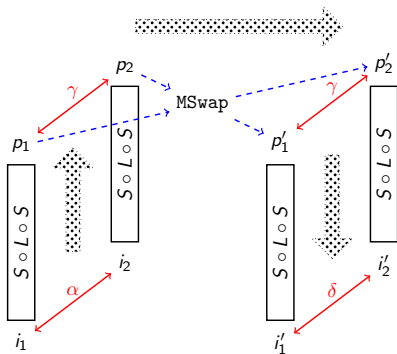
$$\Pr[\exists i : wt(\nu^2(\Delta_i)) = 4] = 2^{-28}$$

## Strategy 2: Composing Improbable/Impossible Differentials

The Inside-Out Technique



Inverted Yoyo

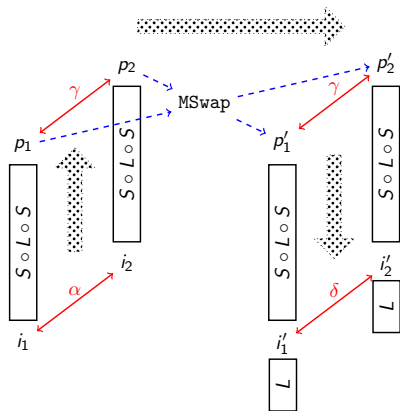


- ▶ By virtue of Yoyo
- ▶  $\Pr[\nu(\alpha) = \nu(\delta)] = 1$

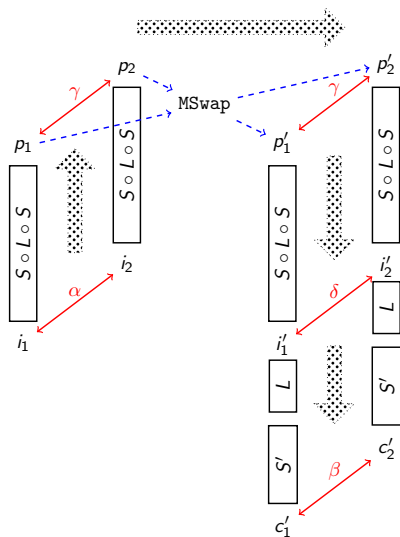
Assumption

Something on  $\nu^2(\delta)$

# Append L



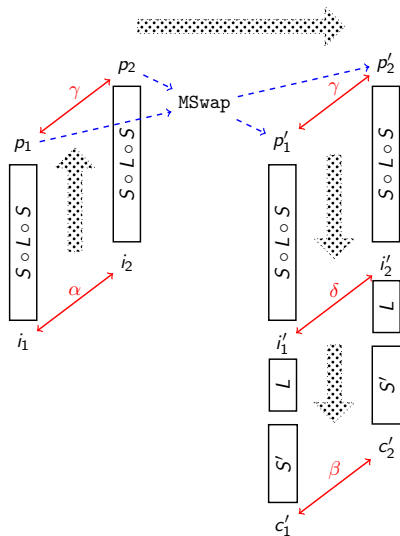
- ▶ Exploit Properties of  $L$
- ▶ Effect of  $L$  on  $\delta$ ?
- ▶ Use  $\nu^2(\delta)$  Assumption



Impossibility

$$\blacktriangleright \Pr[\nu^2(\delta) \rightarrow \nu^2(\beta)] = 0$$

# Probability of $\nu^2(\delta)$ Assumption Holding



## Improbable Differential Yoyo

►  $\Pr[\nu^2(\delta) \text{ Assumption}] < 1$

## Impossible Differential Yoyo

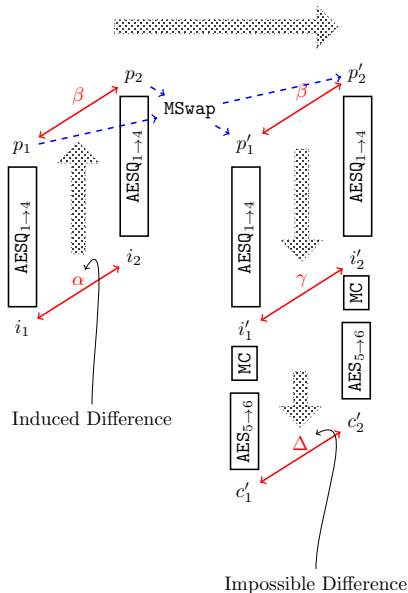
►  $\Pr[\nu^2(\delta) \text{ Assumption}] = 1$

## Application: AES, AESQ

6 Round AES (Practical)

9-10(Practical), 12 Round AESQ

# Impossible Differential Yoyo Distinguisher on 6-Round AES

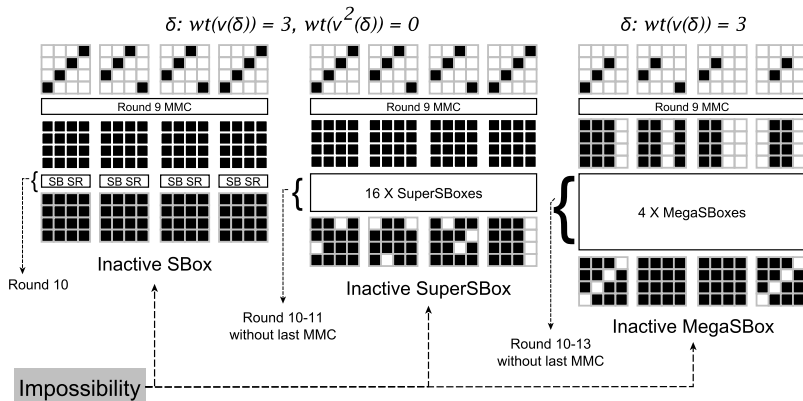


- ▶ One SuperSBox active in  $\alpha$
- ▶ One SuperSBox active in  $\gamma$
- ▶ At least one byte active in  $\gamma$
- ▶ **At least one column active after MC**
- ▶ All SuperSBoxes active after MC

Impossible

One **inactive** SuperSBox in  $\Delta$

## Exploiting Same Property of MixColumns

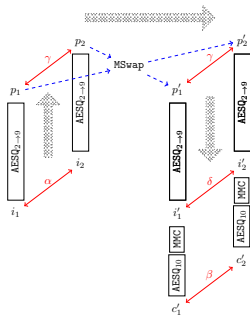


## Impossibilities with different $S'$ Layers

# 9, 10, 11— Round Distinguishers on AESQ

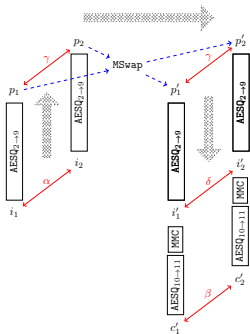
One active word (MegaSBox) in  $\alpha$

## Improbable Differential



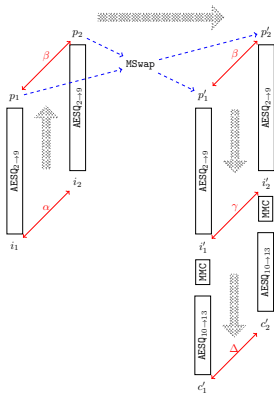
$$\text{AESQ}_{2 \rightarrow 10} : 2^2$$

## Impossible Differential



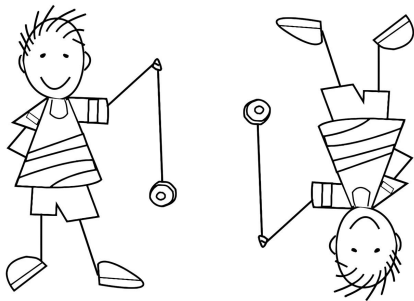
$$\text{AESQ}_{2 \rightarrow 11} : 2^{28}$$

## Impossible Differential



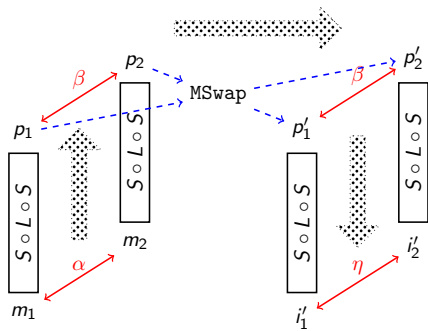
$$\text{AESQ}_{2 \rightarrow 13} : 2^{126}$$

## Strategy 3: Bi-directional Yoyo

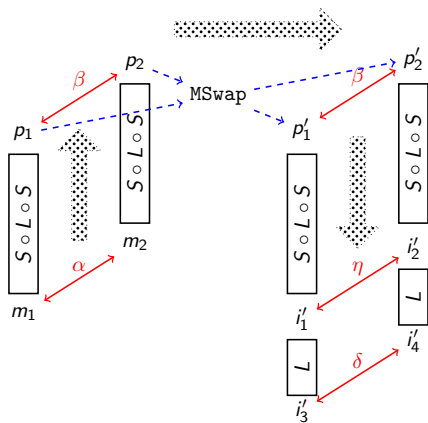


Composing Two Yoyo Games In Two-Directions

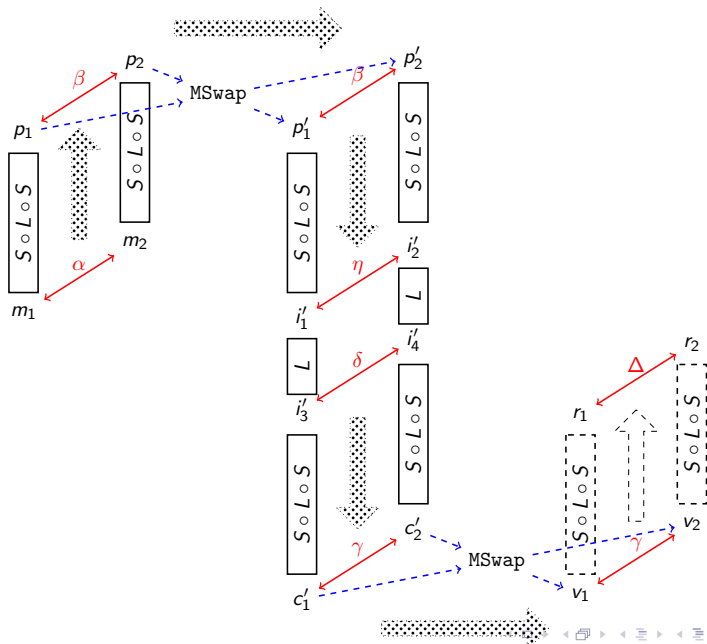
# Inverted Yoyo



# Adding Linear Layer



# Composing 2nd Yoyo

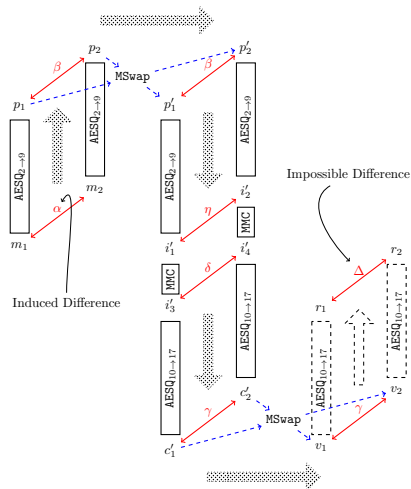
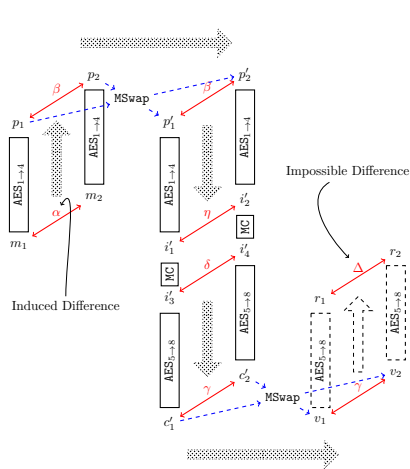




## Application: AES, AESQ

8 Round AES (Practical)

16 Round AESQ



## Distinguishing Complexities

$$\text{AES}_{1 \rightarrow 8} : 2^{30}$$

$$\text{AESQ}_{2 \rightarrow 17} : 2^{126}$$

# Distinguishers on AESQ

Rounds	Complexity		Technique	Reference	
	Time	Memory			
8	$2^{32}$		CICO	Designers	
$8^\dagger$	1	Negligible	YoYo	<b>This Work</b>	
9	$2^{26}$	Negligible		<b>This Work</b>	
$9^\dagger$	$2^2$	Negligible	<b>Improbable Differential YoYo</b>	<b>This Work</b>	
$10^\dagger$	$2^{28}$	Negligible		<b>This Work</b>	
$12^\dagger$	$2^{126}$	Negligible	<b>Impossible Differential YoYo</b>	<b>This Work</b>	
	$2^{256}$	$2^{256}$		Rebound Attack	Designers
	$2^{128}$	Negligible	Time-memory Trade-off		Bagheri et al.
	$2^{102.4}$	$2^{102.4}$			
	$2^{128-x/4}$	$2^x$			
$16^\dagger$	$2^{192}$	$2^{128}$	Rebound Attack		
	$2^{188}$	$2^{128}$	Multi Ltd.-Birthday Distinguisher		
	$2^{192+x}$	$2^{128-x}$	Time-memory Trade-off		
	$2^{126}$	Negligible	<b>Impossible Differential Bidirectional YoYo</b>	<b>This Work</b>	

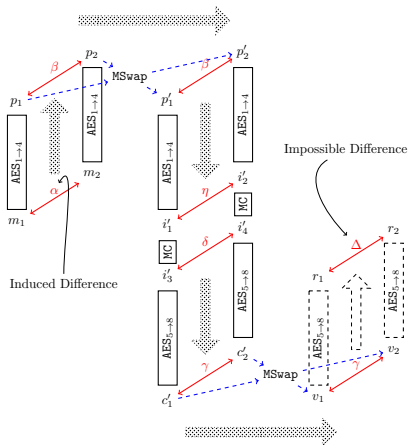
# 8-round Known-Key Distinguishers on AES

Time Complexity	Memory Complexity	Property	Reference
$2^{64}$	$2^{64}$	Uniform Distribution	Gilbert et al.
$2^{48}$	$2^{32}$	Differential Trail	Gilbert et al.
$2^{44}$	$2^{32}$	Multiple Differential Trail	Jean et al.
$2^{30}$	<b>negligible</b>	<b>Impossible Differential Bi-directional Yoyo</b>	<b>This Work</b>
$2^{23}$	$2^{16}$	Extended 7-Round Multiple Differential Trail	Grassi et al.

# Distinguishers reported in this work

	#R	Start → End	Complexity	Strategy	Remarks
AESQ	8	2 → 9	1	Yoyo	Basic Yoyo
	9	1 → 9	$2^{26}$	Yoyo + Nested ZDP	First 9 round Distinguisher starting from Round 1
	9	2 → 10	$2^2$	Improbable Differential Yoyo	Uses the inside-out technique
	10	2 → 11	$2^{28}$		
	12	2 → 13	$2^{126}$	Impossible Differential Yoyo	
	16	2 → 17	$2^{126}$	Bi-directional Impossible Differential Yoyo	Uses inside-out with bi-directional Yoyo
AES	6	1 → 6	$2^{30}$	Impossible Differential Yoyo	Uses the inside-out technique
	8	1 → 8	$2^{30}$	Bi-directional Impossible Differential Yoyo	Uses inside-out with bi-directional Yoyo

# Significance of New Distinguishers



- ▶ Non-triviality
- ▶ Randomness in Queries

- ▶ New ways to extend basic Yoyo game
  - ▶ Classical Differentials
  - ▶ Improbable/Impossible Differentials
  - ▶ Bi-directional Yoyo
- ▶ Using public permutations
- ▶ Best results achieved for AESQ
- ▶ New known-key distinguishers for AES
- ▶ All practical distinguishers experimentally verified
- ▶ Yoyo seems to be an effective generic cryptanalysis tool

**“Almost”** Accepted at IACR ToSC (FSE 2019)



Danke Subria \* TAKK \* Merci  
Xie Xie! THANK YOU  
EFHARISTO TODA  
grazi \* Tack SHUKRAN  
GRACIAS KIITOS  
em INSTUTIIY no Dakkil